# ISO/IEC DTR 13211–1:2006 New built-in flags, predicates, and functions proposal

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October 5, 2009

# Introduction

This proposal specifies a set of built-in predicates and flags to be added to Part 1 of the International Standard for Prolog, ISO/IEC 13211. When evaluating this proposal, please comment each predicate individually by presenting your arguments for either accepting or rejecting its inclusion in the next revision of the Part 1 standard.

This proposal is written as an extension to the ISO/IEC 13211–1 Prolog standard, adopting a similar structure. Specifically, this proposal either adds new sections and clauses to, or modifies the reading of existing clauses on ISO/IEC 13211–1.

This draft proposal may contain in several places informative text, type-set in *italics*. Such informative text is used for editorial comments deemed useful during the development of this draft and may not be included in the final version.

# Contributors

This list includes so far the people present at the WG17 ISO meetings collocated with the ICLP'06 and ICLP'09 conferences and people participating on the mailing list discussions.

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# 1 Scope

This proposal is designed to promote the applicability and portability of Prolog by adding to ISO/IEC 13211–1:1995 a set of built-in predicates and flags that are either common practice and implemented in most Prolog systems or are needed to clarify implementation-dependent behavior. As such, this proposal includes specifications for:

- a) A set of flags allowing a programmer to query a system about the floatingpoint arithmetic implementation and to declare the default encoding for Prolog text
- b) Commonly used term testing predicates already available in most Prolog systems that are missing from ISO/IEC 13211–1:1995
- c) Commonly used meta-predicates which should be available as built-in predicates in order to provide adequate performance
- d) A set of built-in predicates for list processing, providing functionality similar to the atomic processing built-in predicates present on ISO/IEC 13211– 1:1995

# 3 DEFINITIONS

e) Commonly used evaluable functors that are missing from ISO/IEC 13211– 1:1995

NOTE — This part of ISO/IEC 13211 will eventually merge with ISO/IEC 13211–1:1995 resulting in a new version of the Part 1 standard.

# 3 Definitions

**3.34** closure: An atom or a compound term.

# 6 Syntax

#### 6.4 The operator table

The infix operator '><' (bitwise exclusive or) is added to the operator table with the same specification as the bitwise or and and bitwise and operators:

Priority Specifier Operators(s) 500  $yfx + - / \setminus / ><$ 

The prefix operator '+' (unary plus) is added to the operator table with the same specification as the unary minus or the bitwise complement operators:

Priority Specifier Operators(s) 200 fy  $+ - \setminus$ 

The infix operator '^' is used in arithmetic expressions to represent integer exponentiation.

# 7 Language concepts and semantics

# 7.1 Types

7.1.6 Related terms

# 7.9 Pair

P is a pair if it is a compound term '-'(Key, Value) where Key and Value are terms.

NOTE — In Prolog text and this part of ISO/IEC 13211 a pair '-'(Key, Value) is normally written as Key-Value or (Key)-(Value) depending on whether or not Key and Value are operators.

# 7.4 Prolog text

# 7.4.2 Directives

7.10 encoding/1

A directive encoding(E) specifies that the Prolog text being prepared for execution uses the encoding E. When used, this directive shall be the first term, on the first line, in a Prolog text with no extra layout characters (6.5.4) before the :-/1 directive operator. Moreover, a single layout character shall be used between the directive operator and the directive functor. No layout characters or comments shall appear between the directive opening and closing parenthesis.

# 7.11 if/1, elif/1, else/0, endif/0

The directives if (Goal), elif(Goal), else, and endif support the conditional of source code.

# 7.8 Control constructs

#### 7.8.11 call/2-N

# 7.1 Description

call(Closure, Arg1, ...) is true iff call(Goal) is true where Goal is constructed by appending Arg1, ... additional arguments to the arguments (if any) of Closure.

#### 7.2 Template and modes

```
call(@callable_term, ?term, ...)
```

# 7.3 Errors

- a) Closure is a variable — instantiation\_error
- b) Closure is neither a variable nor a callable term
   type\_error(callable, Closure)
- c) The number of arguments in the resulting goal exceeds the implementation defined maximum arity for compound terms
   representation\_error(max\_arity)

# 7.4 Examples

# 7 LANGUAGE CONCEPTS AND SEMANTICS

System	N	Notes
B-Prolog	10/65535	(interpreter/compiler i.e. maximum arity)
CxProlog	9	- · · · ·
GNU Prolog	11	-
SICStus Prolog	255	(maximum arity)
SWI- $Prolog$	8	(meta_predicate/1 directive limited to predicates of arity up to 8)
XSB	11	-
$Y\!AP$	12	-

Table 1: Status of current Prolog implementations providing built-in support for the call/2-N control construct.

```
call(integer, 3).
   Succeeds.
call(atom_concat(pro), log, Atom).
   Succeeds, unifying Atom with prolog.
```

# 7.8.12 setup\_call\_cleanup/3

For details on the specification proposal for this built-in predicate please check: http://www.complang.tuwien.ac.at/ulrich/iso-prolog/cleanup

# 7.10 Input/ouput

# 7.10.3 Standard Streams

Added user\_error alias.

Three streams are predefined and open during the execution of every goal: the standard input stream has the alias user\_input, the standard output stream has the alias user\_output, and the standard error output stream has the alias user\_error. NOTES

3 Prolog implementation on systems that do not support an error output stream shall redirect output to the standard output stream.

# 7.11 Options on stream creation

#### 7.11 Options on stream creation

An implementation may optionally support the following stream-options:

bom(Bool) — If Bool (7.1.4.2) is true then a Unicode encoding Byte Order Mark shall be written when opening a text stream for writing in mode write or is probed for when opening the text stream for reading. This option shall be ignored when opening a stream in mode append. If Bool is false then a Unicode encoding Byte Order Mark shall not be written when opening the text stream for writing.

When no bom(Bool) stream-option is specified, the default value shall be true when the text stream is opened for reading and false when the text stream is opened for writing.

encoding(Encoding) — Encoding is an atom representing the text encoding that shall be used when opening the stream for writing or the text encoding of the stream opened for reading.

When opening the text stream for reading with the default bom(Bool) streamoption value or by explicitly specifying the bom(true) stream-option, if a *Byte Order Mark* is detected, it will be used to set the corresponding Unicode text stream encoding, overriding any encoding(Encoding) that might be also specified.

### NOTES

- 1 These stream-options imply the stream-option type(text).
- 2 The set of supported text encodings is implementation-defined.

# 7.13 Stream properties

An implementation may optionally support the following stream properties:

bom(Bool) — If present and if Bool (7.1.4.2) is true, a Unicode encoding Byte Order Mark was detected while opening the text stream for reading or a Byte Order Mark was written while opening the text stream for writing.

encoding(Encoding) — Encoding used for the text stream.

NOTE — These stream properties imply the stream property type(text).

# 7.11 Flags

7.11.2 Flags defining float type F

7.1 Flag: float\_mantissa\_digits

# 7 LANGUAGE CONCEPTS AND SEMANTICS

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is the number of significant digits on the mantissa of a normalized floating point number (in base 10), an implementation defined integer value.

# 7.2 Flag: float\_epsilon

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is the distance from 1.0 to the next largest floating point number, an implementation defined value. Thus, it allows the programmer to query an implementation about the relative accuracy when performing arithmetic with floating point numbers.

#### 7.3 Flag: float\_min\_exponent

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is smallest possible value of the exponent of a normalized floating point number, an implementation defined integer value.

# 7.4 Flag: float\_max\_exponent

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is greatest value of the exponent of a normalized floating point number, an implementation defined integer value.

# 7.11.3 Other flags

#### 7.6 Flag: unification\_subject\_to\_occurs\_check

Possible values: fail, cyclic, unsafe

Default: implementation defined

Changeable: implementation defined

Description: This read-only flag describes the behaviour of the Prolog system when a variable is unified with a compound term that contains it (STO unification, 3.165). The flag value fail implies that STO unification simply fails. The flag value unsafe means that if an STO unification is encountered the further behavior of the system is undefined. The flag value cyclic implies that STO unifications will be successful and result in the creation of cyclic terms. Moreover, this flag indicates that unification between cyclic terms terminates.

#### NOTE

The flag value fail means that the built-in predicate =/2 (8.2.1) behaves exactly as the built-in predicate unify\_with\_occurs\_check/2 (8.2.2). The flag value unsafe may imply that the STO unification itself, or further unifications or built-in predicate calls may not terminate, or cause the system to fail or raise an exception.

Examples:

| ?- X = f(X).

- fail: fails
- cyclic: succeeds and unifies X with a cyclic term f(f(f(...))).
- $\bullet$  unsafe: undefined. Often succeeds, but subsequent use of X, as e.g. in X=X, causes an error.

| ?- X = f(X), Y = f(Y), X = Y. | ?- g(X,Y,X) = g(f(X),f(Y),Y). | ?- X = f(X), Y = f(Y), X == Y. | ?- X = f(X), asserta(p(X)).

For all the above four examples:

- fail: fails
- cyclic: succeeds and unifies both X and Y with a cyclic term f(f(f(...))).
- unsafe: undefined. Often causes an error.

# 7.7 Flag: encoding

Possible values: an implementation defined atom

Default: implementation defined

Changeable: implementation defined

Description: This flag represents the default encoding for text streams. An implementation shall document if the flag value can be changed by programmer as well all the supported encodings. The text encoding names to be used shall be the ones specified by the Internet Assigned Numbers Authority (IANA) and marked as "(preferred MIME name)":

http://www.iana.org/assignments/character-sets

NOTE — Implementations can define aliases for the standard text encodings names.

# 7.12 Errors

# 7.12.2 Error classification

The following types are added to the classification of 7.12.2 of ISO/IEC 13211-1.

- a) The list of valid types is extended by the addition of pair (see 7.12.2 b of ISO/IEC 13211-1).
- b) The list of valid domains is extended by the addition of order and predicate\_property (see 7.12.2 c of ISO/IEC 13211-1).

# 7.13 Predicate properties

The properties of procedures can be found using the built-in predicate predicate\_property(Callable, Property), where Callable is a callable term. The predicate properties supported shall include:

- **static** The predicate is static
- dynamic The predicate is dynamic
- built\_in The predicate is a built-in predicate
- multifile The predicate is the subject of a multifile directive

A processor may support one or more additional predicate properties as an implementation specific feature. Implementation-defined properties are not required to be atomic terms.

# 8 Built-in predicates

The following sections extends, with the specified number, the corresponding ISO/IEC 13211–1 sections:

# 8.2 Term unification

#### 8.2.4 subsumes/2

# 8.1 Description

subsumes (General, Specific) is true iff there is a substitution  $\theta$ , including the empty substitution, such that the term General is instantiated to  $General\theta = Specific\theta$ . This predicate provides a one-way unification.

### 8.2 Template and modes

subsumes(?term, @term)

#### 8.3 Errors

None.

# 8.4 Examples

```
subsumes(f(X,Y), f(Z,Z)).
Succeeds, unifying both X and Y to Z.
```

subsumes(f(Z,Z), f(X,Y)).
Fails.

# 8.3 Type testing

8.3.9 callable/1

# 8.1 Description

callable(Term) is true iff Term is a callable term.

#### 8.2 Template and modes

callable(@term)

#### 8.3 Errors

None.

# 8.4 Examples

callable(a). Succeeds.

callable(3). Fails.

callable(X).
Fails.

8.3.10 ground/1

# 8.1 Description

ground(Term) is true iff Term is a ground term.

# 8.2 Template and modes

ground(@term)

# 8.3 Errors

None.

# 8.4 Examples

ground(3). Succeeds.

ground(a(1, \_)).
Fails.

# 8.3.11 acyclic\_term/1

# 8.1 Description

acyclic\_term(Term) is true iff Term is an acyclic term. For implementations not supporting STO unification 7.6, calls to this predicate simply succeed.

# 8.2 Template and modes

acyclic\_term(@term)

# 8.3 Errors

None.

#### 8.4 Examples

acyclic\_term(a(1, \_)).
Succeeds.

X = f(X), acyclic\_term(X).
Fails.

8.3.12 cyclic\_term/1

### 8.1 Description

cyclic\_term(Term) is true iff Term is a cyclic term. For implementations not supporting STO unification 7.6, calls to this predicate simply fail.

#### 8.2 Template and modes

cyclic\_term(@term)

#### 8.3 Errors

None.

# 8.4 Examples

cyclic\_term(a(1, \_)).
Fails.

X = f(X), cyclic\_term(X). Succeeds.

# 8.4 Term comparison

8.4.2 compare/3

### 8.1 Description

compare(Order, Term1, Term2) is true iff Order corresponds to the standard order between Term1 and Term2. The argument Order is unified with the atom < when Term1 is less than Term2, with the atom = when Term1 and Term2 are equal, and with the atom > when Term1 is greater than Term2.

#### 8.2 Template and modes

compare(?atom, @term, @term)

#### 8.3 Errors

- a) Order is neither a variable nor an atom
   type\_error(atom, Order)

#### 8.4 Examples

```
compare(Order, 3, 5).
Succeeds, unifying Order with <.
compare(Order, d, d).
Succeeds, unifying Order with =.
compare(Order, 3, 3.0).
Succeeds, unifying Order with >.
```

# 8.5 Term creation and decomposition

```
8.5.5 numbervars/3
```

8.1 Description

numbervars(Term, Start, End) is true. This predicate unifies each free variable on Term with a compound term with the format '\$VAR'(N) where N is an integer starting from Start and ending at End-1.

#### 8.2 Template and modes

numbervars(?nonvar, +integer, -integer)

#### 8.3 Errors

- a) Start is a variable — instantiation\_error
- b) Start is neither a variable nor an integer— type\_error(integer, Start)

#### 8.4 Examples

```
numbervars(foo(A, B, A), 0, End).
Succeeds, unifying A with '$VAR'(0), B with '$VAR'(1),
and End with 2.
```

# 8.8 Clause retrieval and information

#### 8.8.3 predicate\_property/2

#### 8.1 Description

predicate\_property(Head, Property) is true iff the procedure associated with the argument Head (3.84) has predicate property Property.

Procedurally, predicate\_property(Head, Property) is executed as follows

- a) Determines the principal functor P and arity N associated with Head. P/N is the associated predicate indicator
- b) Searches the complete database and creates a set SetPP of all terms PP such that P/N identifies a procedure which has predicate property PP and PP is unifiable with Property
- c) If SetPP is non empty set proceeds to 8.1 e,
- d) Else the goal fails
- e) Chooses the first element PPP of *SetPP*, unifies PPP with Property and the predicate succeeds
- f) If all elements of SetPP have been chosen the predicate fails
- g) Else chooses the first element PPP of *SetPP* that has not already been chosen, unifies PPP with Property and the goal succeeds

predicate\_property(Head, Property) is re-executable. On backtracking, continue at 8.1 f.

The order in which properties are found by predicate\_property/2 is implementation dependent.

#### NOTES

1 A processor may support, as an implementation specific feature, additional predicate properties.

2 For a dynamic predicate, all proprieties related to its definition shall be removed when the predicate is abolished.

#### 8.2 Template and modes

predicate\_property(?callable\_term, ?predicate\_property)

- 8.3 Errors
  - a) Head is neither a variable nor a callable term type\_error(callable, Head)
  - b) Property is neither a variable nor a predicate property
     domain\_error(predicate\_property, Property)

#### 8.4 Examples

```
predicate_property(once(_), built_in).
   Succeeds.
predicate_property(atom_codes(_, _), Property).
   Succeeds unifying Property with static.
   On re-execution, succeeds unifying Property with built_in.
```

# 8.9 Clause creation and destruction

# 8.9.3 retract/1

# 8.1 Errors

There is a typo on the current standard in the specification of the permission\_error exception that should use the atom modify instead of access in order to match the specification of other database predicates.

This typo may or may not be already corrected in the 13211-1 errata recently published.

- a) ...
- b) ...
- c) The predicate indicator Pred of Head is that of a static procedure — permission\_error(modify, static\_procedure, Pred)

# 8.9.5 retractall/1

# 8.1 Description

retractall(Head) is true. Procedurally, the predicate shall behave as if defined by:

```
retractall(Head) :-
    retract((Head :- _)),
    fail.
retractall(_).
```

Procedurally, retractall(Head) is executed as follows:

- a) Determines the principal functor P and arity N associated with Head. P/N is the associated predicate indicator
- b) If the database contains a dynamic procedure whose predicate indicator is P/N, then proceeds to 8.1 d,
- c) Else the goal succeeds.
- d) Retracts from the database all clauses whose head unifies with Head and the goal succeeds

# NOTES

1 The dynamic predicate shall continue to be known by the system even when all of its clauses are removed.

2 This predicate does not change any of the standard predicate properties of the referenced predicate (as reported by predicate\_property(Head, Property)), even when all of its clauses are removed.

#### 8.2 Template and modes

retractall(@callable\_term)

# 8.3 Errors

- a) Head is a variable — instantiation\_error
- b) Head is neither a variable nor a callable term
   type\_error(callable, Generate)
- c) The predicate indicator Pred of Head is that of a static procedure
   permission\_error(modify, static\_procedure, Pred)

#### 8.4 Examples

The examples defined in this subclause assume the database has been created from the following Prolog text:

```
:- dynamic(insect/1).
insect(ant).
insect(bee).
insect(spider).
retractall(insect(bee)).
Succeeds, retracting the clause 'insect(bee)'.
retractall(insect(_)).
Succeeds, retracting all the clauses of predicate insect/1.
retractall(insect(elephant)).
Succeeds.
retractall(mammal(_)).
Succeeds.
retractall(3).
type_error(callable, 3)
```

# 8.10 All solutions

8.10.4 forall/2

8.10.4.1 Description

forall(Generate, Test) is true iff for all possible bindings of Generate, the goal Test is true. Procedurally, the predicate shall behave as if defined by \+ (call(Generator), \+ call(Test)).

#### 8.10.4.2 Template and modes

forall(@callable\_term, @callable\_term)

#### 8.10.4.3 Errors

- a) Generate is a variable — instantiation\_error
- b) Generate is neither a variable nor a callable term
   type\_error(callable, Generate)
- c) Test is a variable — instantiation\_error
- d) Test is neither a variable nor a callable term
   type\_error(callable, Test)

#### 8.10.4.4 Examples

The following examples assume that the predicate a/1 and b/1 are defined with the following clauses:

```
a(1). a(2). a(3).
b(1, a). b(2, b). b(3, c).
forall(fail, true).
Succeeds.
forall(a(X), b(X, _)).
Succeeds.
forall(a(X), b(_, X)).
Fails.
forall(b(_, Y), write(Y))
Succeeds, outputting the characters
abc
to the current output stream.
```

#### 8.10.5 sort/2

#### 8.10.5.1 Description

sort(List, Sorted) is true iff Sorted is a list containing the non-duplicated elements of List sorted in ascending order following standard order (7.2).

#### 8.10.5.2 Template and modes

sort(@list, ?list)

#### 8.10.5.3 Errors

- a) List is a partial list — instantiation\_error
- b) List is neither a partial list nor a list
   type\_error(list, List)
- c) Sorted is neither a partial list nor a list
   type\_error(list, Sorted)

# 8.10.5.4 Examples

```
sort([1, 2, 1, 8, 4], Sorted).
Succeeds, unifying Sorted with [1, 2, 4, 8].
```

#### 9 EVALUABLE FUNCTORS

#### 8.10.6 keysort/2

#### 8.10.6.1 Description

keysort(List, Sorted) is true iff List is a list of elements with the format Key-Value and Sorted is a list containing the elements of List sorted according to the value of Key in ascending order following standard order (7.2). The sorting procedure shall be stable. I.e. the relative order of elements of List with the same key shall not change in the Sorted list.

#### 8.10.6.2 Template and modes

keysort(@list, ?list)

#### 8.10.6.3 Errors

- a) List is a partial list — instantiation\_error
- b) List is neither a partial list nor a list
   type\_error(list, List)
- c) An element Element of List is a variable instantiation\_error
- d) An element  $\tt Element$  of  $\tt List$  is neither a variable nor a '-'/2 compound term

— type\_error(pair, Element)

e) Sorted is neither a partial list nor a list
 — type\_error(list, Sorted)

# 8.10.6.4 Examples

```
keysort([1-a, 3-f(_), 1-z, 2-44], Sorted).
Succeeds unifying Sorted with [1-a, 1-z, 2-44, 3-f(_)].
```

# 9 Evaluable functors

# 9.1 The simple arithmetic functors

The unary plus evaluable arithmetic functor is added.

# 9.1.1 Evaluable functors and operations

Evaluable functor Operation

(+)/1  $pos_I, pos_F$ 

# 9.1.3 Integer operations and axioms

The following operations are specified:

 $pos_I: I \to I$ 

For all  $x \in I$ , the following axioms shall apply:

 $pos_I(x) = x$ 

# 9.1.4 Floating point operations and axioms

The following operations are specified:

 $pos_F: F \to F$ 

For all  $x \in F$ , the following axioms shall apply:

 $pos_F(x) = x$ 

# 9.3 Other arithmetic functors

 $9.3.8 ^{/2}$ 

# 9.3.8.1 Description

 $^(VI, VJ)$  evaluates the expression I with value VI, the expression J with value VJ, and has the value of the integer exponentiation of VI to VJ.

#### 9.3.8.2 Template and modes

^(int-exp, int-exp) = integer

#### 9.3.8.3 Errors

- a) I is a variable — instantiation\_error
- b) I is not a variable and VI is not an integer
   type\_error(integer, VI)
- c) J is a variable -- instantiation\_error
- d) J is not a variable and VJ is not an integer
   type\_error(integer, VJ)

# 9.3.8.4 Examples

34<sup>12</sup>. Evaluates to 2386420683693101056.

9.3.9  $\log/2$ 

#### 9.3.9.1 Description

log(B, X) evaluates the expression B with value VB, the expression X with value VX, and has the value of the logarithm to base VB of VX.

#### 9.3.9.2 Template and modes

log(int-exp, float-exp) = float
log(int-exp, int-exp) = float

### 9.3.9.3 Errors

- a) B is a variable — instantiation\_error
- b) B is not a variable and VB is not an integer— type\_error(integer, VB)
- c) VB is zero or negative
   evaluation\_error(undefined)
- d) X is a variable — instantiation\_error
- e) X is not a variable and VX is not a number
   type\_error(number, VX)

# 9.3.9.4 Examples

log(10, 10.0). Evaluates to 1.0.

#### $9.3.10 \ \mathrm{gcd}/2$

# 9.3.10.1 Description

gcd(I, J) evaluates the expression I with value VI, the expression J with value VJ, and has the value of the greatest common divisor of VI of VJ.

# 9 EVALUABLE FUNCTORS

#### 9.3.10.2 Template and modes

gcd(int-exp, int-exp) = integer

#### 9.3.10.3 Errors

- a) I is a variable instantiation\_error
- b) I is not a variable and VI is not an integer
   type\_error(integer, VI)
- c) J is a variable instantiation\_error
- d) J is not a variable and VJ is not an integer
   type\_error(integer, VJ)

#### 9.3.10.4 Examples

gcd(2, 3). Evaluates to 1.

# $9.3.11 \mod 2$

#### 9.3.11.1 Description

 $\max(X, Y)$  evaluates the expression X with value VX, the expression Y with value VY, and has the value of the maximum of VX and VY. When used with expressions of mixed-types, the result is implementation-dependent; an implementation may chose either to return a value or to throw an exception.

#### 9.3.11.2 Template and modes

```
max(float-exp, float-exp) = float
max(float-exp, int-exp) = implementation-dependent result
max(int-exp, float-exp) = implementation-dependent result
max(int-exp, int-exp) = integer
```

#### 9.3.11.3 Errors

- a) X is a variable or Y is a variable
   instantiation\_error
- b) X is not a variable and VX is not a number
   type\_error(number, VX)
- c) Y is not a variable and VY is not a number— type\_error(number, VY)

9.3.11.4 Examples

```
max(2, 3)
    Evaluates to 3.
max(2.0, 3.0)
    Evaluates to 3.0.
max(0, 0.0).
```

Implementation-dependent result.

#### 9.3.12 min/2

#### 9.3.12.1 Description

min(X, Y) evaluates the expression X with value VX, the expression Y with value VY, and has the value of the minimum of VX and VY. When used with expressions of mixed-types, the result is implementation-dependent; an implementation may chose either to return a value or to throw an exception.

#### 9.3.12.2 Template and modes

```
min(float-exp, float-exp) = float
min(float-exp, int-exp) = implementation-dependent result
min(int-exp, float-exp) = implementation-dependent result
min(int-exp, int-exp) = integer
```

# 9.3.12.3 Errors

- a) X is a variable or Y is a variable
   instantiation\_error
- b) X is not a variable and VX is not a number
   type\_error(number, VX)
- c) Y is not a variable and VY is not a number— type\_error(number, VY)

# 9.3.12.4 Examples

```
min(2, 3)
    Evaluates to 2.
min(2.0, 3.0)
    Evaluates to 2.0.
```

min(0, 0.0).
Implementation-dependent result.

# 9.4 Bitwise functors

9.4.6 (><)/2 – bitwise exclusive or

#### 9.4.6.1 Description

 $^{\prime}$  (B1, B2) evaluates the expressions B1 and B2 with values VB1 and VB2 and has the value such that each bit is set iff only one of the corresponding bits in VB1 and VB2 is set.

The value shall be implementation defined if VB1 or VB2 is negative.

#### 9.4.6.2 Template and modes

'><'(int-exp, int-exp) = integer</pre>

NOTE — '><' is an infix predefined operator (see 6.3.4.4).

# 9.4.6.3 Errors

- a) B1 is a variable — instantiation\_error
- b) B2 is a variable — instantiation\_error
- c) B1 is not a variable and VB1 is not an integer— type\_error(integer, VB1)
- d) B2 is not a variable and VB2 is not an integer
   type\_error(integer, VB2)

# 9.4.6.4 Examples

```
'><'(10, 12).
Evaluates to the value 6.
'><'(125, 255).
Evaluates to to the value 130.
'><'(-10, 12).
Evaluates to an implementation defined value.
'><'(77, N)
instantiation_error.
'><'(foo, 2)
type_error(integer, foo).
```

# 9.5 Trigonometric functors

Assuming that we will be adding the trigonometric functions described below to the revised core standard, it's probably best if we gather all the trigonometric functions under their own section. The following order is assumed below:  $\sin/1$ ,  $\cos/1$ ,  $\tan/1$ ,  $a\sin/1$ ,  $a\cos/1$ ,  $a\tan/2$ .

# $9.5.3 \tan/1$

# 9.5.3.1 Description

tan(X) evaluates the expression X with value VX and has the value of the tangent of VX (measured in radians).

#### 9.5.3.2 Template and modes

```
tan(float-exp) = float
tan(int-exp) = float
```

#### 9.5.3.3 Errors

- a) X is a variable -- instantiation\_error
- b) X is not a variable and VX is not a number
   type\_error(number, VX)

# 9.5.3.4 Examples

tan(pi).
Evaluates to 0.0.

# $9.5.4 \ asin/1$

#### 9.5.4.1 Description

asin(X) evaluates the expression X with value VX and has the value of the arc sine of VX (in radians).

#### 9.5.4.2 Template and modes

asin(float-exp) = float
asin(int-exp) = float

# 9.5.4.3 Errors

- a) X is a variable — instantiation\_error
- b) X is not a variable and VX is not a number
   type\_error(number, VX)

# 9.5.4.4 Examples

asin(1.0).

Evaluates to a value approximately equal to 1.570796326795.

 $9.5.5 \ acos/1$ 

#### 9.5.5.1 Description

acos(X) evaluates the expression X with value VX and has the value of the arc cosine of VX (in radians).

### 9.5.5.2 Template and modes

```
acos(float-exp) = float
acos(int-exp) = float
```

#### 9.5.5.3 Errors

- a) X is a variable — instantiation\_error
- b) X is not a variable and VX is not a number— type\_error(number, VX)

# 9.5.5.4 Examples

acos(0.0). Evaluates to a value approximately equal to 1.570796326795.

#### $9.5.6 ext{ atan}/2$

# 9.5.6.1 Description

atan(Y, X) evaluates the expression Y with value VY, the expression X with value VX, and computes the principal value of the arc tangent of VY/VX (in radians), using the signs of both arguments to determine the quadrant of the return value. When both arguments are 0.0, an implementation-dependent floating-point value is returned.

# 9 EVALUABLE FUNCTORS

# 9.5.6.2 Template and modes

```
atan(float-exp, float-exp) = float
atan(float-exp, int-exp) = float
atan(int-exp, float-exp) = float
atan(int-exp, int-exp) = float
```

# 9.5.6.3 Errors

- a) X is a variable — instantiation\_error
- b) Y is not a variable and VY is not a number— type\_error(number, VY)
- c) X is not a variable and VX is not a number
   type\_error(number, VX)

#### 9.5.6.4 Examples

```
atan(0.0, -0.0).
Evaluates to a value approximately equal to 3.14159265358979.
```

# 9.6 Mathematical constants

# 9.6.1 pi/0

#### 9.6.1.1 Description

pi evaluates to the floating-point number which best approximates the mathematical constant  $\pi$ , the ratio of a circle's circumference to its diameter.

#### 9.6.1.2 Examples

# pi.

Evaluates to the corresponding mathematical constant. The accuracy of the result is implementation-defined.

#### 9.6.2 e/0

# 9.6.2.1 Description

**e** evaluates to the floating-point number which best approximates the mathematical constant *e*, the base of natural logarithms.

# 9.6.2.2 Examples

e.

Evaluates to the corresponding mathematical constant. The accuracy of the result is implementation-defined.

9.6.3 epsilon/0

# 9.6.3.1 Description

**epsilon** evaluates to the distance from 1.0 to the next largest floating point number, an implementation defined value. Thus, it allows the programmer to retrieve the relative accuracy when performing arithmetic with floating point numbers.

# 9.6.3.2 Examples

epsilon. Evaluates to an implementation defined value.